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A Measure of Robustness Against Multiple Kinds of Perturbations

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Abstract

Parallel and distributed heterogeneous computing systems may operate in an environment that undergoes unpredictable changes causing certain system performance features to degrade. Such systems need robustness to guarantee limited degradation despite fluctuations in the behavior of its component parts or environment. Our previous work in this area presented a method for generating a measure of robustness for a given system. However, the focus of that approach was on a scenario where all perturbations were of the same kind, e.g., all perturbations were in message sizes or computation times, but not both message sizes and computation times. This paper gives an extended discussion of the case where perturbations could be of different kinds, and presents some new insights.

Keywords: robustness, robustness metric, resource allocation, resource management systems, parallel and distributed systems.

1. Introduction

The robust design of computing and communication systems is becoming an increasingly important issue [1, 3–15]. There is a need for research that addresses the issue of developing a generalized robustness metric. In this paper, we extend our previous formulation ([2]) of a standard generalized robustness metric for resource allocation. This will be an important step towards ongoing efforts to create robust designs.

The motivation for this research was provided by research supported by the DARPA's ITO Quorum program, under the project called "Management System for Heterogeneous Networks." The research involved the design and analysis of heuristics for robust resource allocation in different types of heterogeneous computing environments including the HiPer-D (High Performance Distributed Computing Program). A typical HiPer-D computing system consists of a set of dedicated machines interconnected by high-speed communication links. A set of sensors (radars, sonars, etc.) sends streams of data sets to a set of communicating, continuously running applications that process these data sets and send their outputs to other applications or actuators.

A HiPer-D system is required to satisfy a set of throughput and latency constraints. Any allocation of the resources must enforce these quality of service (*QoS*) constraints by ensuring that the computation and communication times are within certain limits. When the system is first configured, it is assumed to operate under certain estimated values of the initial sensor loads (i.e., outputs from sensors). Such an initial resource allocation ensures that all throughput and latency constraints are met when the system is first deployed. However, the system is expected to operate in a dynamic environment, where the sensor loads are expected to change unpredictably. Increases in sensor loads cause increases in the computation and communication times, which in turn may cause throughput and latency violations. Therefore, the initial resource allocation might be rendered invalid soon after the operation begins.

One way of handling the unpredictable load increases is to design a resource allocation that will tolerate as much sensor load increase as possible before a QoS violation occurs. For such an approach, how does one determine which resource allocation tolerates the largest load increase, given a set of resource allocations? This task necessitates the formulation of an appropriate metric.

One needs a general approach because the sensor loads might not be the only uncertainties in a HiPer-D system. Two other examples are: (a) inaccurate models for computation/communication times, and (b) sudden machine or link failures. A general approach is necessary also because for systems other than HiPer-D, there might be other uncertainties. Typically, the resource allocation decisions and the performance prediction are based on estimated/initial values of application and system parameters. However, complex computing and communication systems typically operate in an unpredictable environment where the actual values of these parameters may differ from the estimates due to a variety of reasons. As a result, the "real" system performance may degrade. An important question then arises. *Given a resource allocation, what is the maximum departure from the expected conditions that the system can tolerate and still deliver the promised performance?* That is, how *robust* is the system? Our research in [2] presented a method for generating a measure of robustness for a given system. However, the focus of that approach was on a scenario where all perturbations were of the same kind, e.g., all perturbations were in message sizes or computation times, but not both message sizes and com-

putation times. The research in [2] did outline a method for generating a robustness measure when the perturbations were of mixed kinds. This paper gives an extended discussion of that method, and presents some new insights.

The rest of the paper is organized as follows. Section 2 summarizes the FePIA procedure given in [2] for generating a robustness measure for an arbitrary system. Section 3 presents our extension of the work in [2] to the scenarios where there are multiple perturbation parameters of different kinds. Section 4 concludes this paper. A glossary of the notation used in this paper is given in Table 1.

2. A Method for Generating Robustness Metrics

The research in [2] proposed a general procedure, called *FePIA,* for deriving a general robustness metric for any desired computing environment. For reference, we summarize the FePIA procedure here. Please see [2] for details.

1) Describe quantitatively the requirement that makes the system robust. Based on this *robustness requirement*, determine the QoS performance features that should be limited in variation to ensure that the robustness requirement is met. Identify the acceptable variation for these feature values as a result of uncertainties in system parameters. Mathematically, let Φ be the set of system performance features that should be limited in variation. For each element $\phi_i \in \Phi$, quantitatively describe the tolerable variation in ϕ_i . Let $\left\langle \beta_i^{\min}, \beta_i^{\max} \right\rangle$ be a tuple that gives the bounds of the tolerable variation in the system feature ϕ_i .

2) Identify all of the system and environment parameters whose values may impact the QoS performance features selected in step 1. These are called the *perturbation parameters*, and the performance features are required to be robust with respect to these perturbation parameters. Mathematically, let Π be the set of perturbation parameters. It is assumed that the elements of Π are vectors. Let π_j be the j-th element of Π . For the makespan example, π_j could be the vector composed of the actual application execution times, i.e., the i -th element of π_i is the actual execution time of the *i*-th application on the machine it was assigned. In general, representation of the perturbation parameters as separate elements of Π would be based on their nature or kind (e.g., message length variables in π_1 and computation time variables in π_2).

3) Identify the impact of the perturbation parameters in step 2 on the system performance features in step 1. Mathematically, for every $\phi_i \in \Phi$, determine the relationship $\phi_i = f_{ij}(\pi_j)$, if any, that relates ϕ_i to π_j . In this expression, f_{ij} is a function that maps π_j to ϕ_i .

4) The last step is to determine the smallest collective variation in the values of perturbation parameters identified in step 2 that will cause any of the performance features identified in step 1 to violate its acceptable variation. This will be the degree of robustness of the given resource allocation.

Mathematically, for every $\phi_i \in \Phi$, determine the *boundary values of* π_j , i.e., the values satisfying the *boundary relationships* $f_{ij}(\pi_j) = \beta_i^{\min}$ and $f_{ij}(\pi_j) =$ β_i^{max} . These relationships separate the region of robust operation from that of non-robust operation. Find the smallest perturbation in π_i that causes any $\phi_i \in \Phi$ to exceed the bounds $\langle \beta_i^{\min}, \beta_i^{\max} \rangle$ imposed on it by the robustness requirement.

Figure 1 illustrates this concept for a single feature, ϕ_i , and a two-element perturbation vector $\pi_j \in \mathbb{R}^2$. The curve shown in Figure 1 plots the set of boundary points $\{\pi_j | f_{ij}(\pi_j) = \beta_i^{\max}\}\$ for a resource allocation μ . For this figure, the set of boundary points $\{\pi_j | f_{ij}(\pi_j) = \beta_i^{\min}\}\$ is given by the points on the π_{j1} axis and π_{i2} -axis.

The region enclosed by the axes and the curve gives the values of π_j for which the system is robust with respect to ϕ_i . For a vector $\mathbf{x} = [x_1 \, x_2 \, \cdots \, x_n]^{\mathrm{T}}$, let $\|\mathbf{x}\|_2$ be the ℓ_2 -norm (Euclidean norm) of the vector, defined by $\sqrt{\sum_{n=1}^{n}}$ $r=1$ x_r^2 . The point on the curve marked as $\pi_j^*(\phi_i)$ has the property that the Euclidean distance from π_j^{orig} to $\pi_j^{\star}(\phi_i)$, $\|\pi_j^{\star}(\phi_i) - \pi_j^{\text{orig}}\|_2$, is the smallest over all such distances from π_j^{orig} to a point on the curve, and is defined to be the *robustness radius*, $r_{\mu}(\phi_i, \pi_j)$, of ϕ_i against π_i . Mathematically,

$$
r_{\mu}(\phi_i, \boldsymbol{\pi}_j) = \min_{\substack{\boldsymbol{\pi}_j: (f_{ij}(\boldsymbol{\pi}_j) = \beta_i^{\text{max}}) \vee (f_{ij}(\boldsymbol{\pi}_j) = \beta_i^{\text{min}})}} \|\boldsymbol{\pi}_j - \boldsymbol{\pi}_j^{\text{orig}}\|_2.
$$
 (1)

The quantity $\rho_{\mu}(\Phi, \pi_i)$ is defined as the robustness of resource allocation μ with respect to the performance feature set Φ against the perturbation parameter π_j , and is given by the minimum of all robustness radii, or $\rho_\mu(\Phi, \pi_j) = \min_{\phi_i \in \Phi} (r_\mu(\phi_i, \pi_j)).$

3. Robustness Against Multiple Perturbation Parameters

3.1. Overview

The research in [2] developed the analysis for determining the robustness metric for a system with a single perturbation parameter. In this section, that analysis is extended to include multiple perturbation parameters.

Consider a HiPer-D like system where task execution times, e_j , and message lengths m_k , change unpredictably. Also assume that the performance feature, ϕ_i , is a function of e_j and m_k . Given such a function for ϕ_i , one could use the method summarized in Section 2 to "determine" $\rho_{\mu}(\Phi, \pi_i)$. However, the method in Section 2 assumes that all elements of π_j have the same units. Because e_i and m_k have different units, one cannot assemble all of them in one perturbation parameter, π_j , and then determine $\rho_\mu(\Phi, \pi_j)$.

Furthermore, [2] makes it clear that the unit of $\rho_{\mu}(\Phi, \pi_i)$ is the same as that for any element of π_i . Given this fact, any calculation for $\rho_{\mu}(\Phi, \pi_j)$ would be questionable even if one did assemble all e_i and m_k in one perturbation parameter. This further clarifies why one cannot assemble e_i and m_k in one π_i without first adjusting for the unit changes.

One can argue that ϕ_i includes necessary conversions of multiple perturbation parameters which are not of the same kind. However, these conversions are not included in the robustness measurement equation. That is the reason why we have to merge π_j 's into **P** so as to compose a dimensionless perturbation parameter vector.

Multiple perturbation parameters are considered in [2] by concatenating them into one parameter, which is then used as a single parameter as discussed in Section 2. Then the robustness metric is determined by taking the minimum over the robustness radii of all $\phi_i \in \Phi$.

Let the vector π_j have n_{π_j} elements, and let \star be the vector concatenation operator, so that $\pi_1 \star \pi_2$ $[\begin{array}{cccc} \pi_{11} & \pi_{12} & \cdots & \pi_{1n_{\pi_1}} & \pi_{21} & \pi_{22} & \cdots & \pi_{2n_{\pi_2}} \end{array}]^{\mathrm{T}}.$ Let $P \in \mathbb{P}$ be a weighted concatenation of the vectors $\pi_1, \pi_2, \cdots, \pi_{|\Pi|}$, where $\mathbb P$ is a space of $n_{\pi_1} + n_{\pi_2} + \cdots + n_{\pi_{|\Pi|}}$ dimensions. That is, $\mathbf{P} = (\alpha_1 \times \pi_1) \star (\alpha_2 \times \pi_2) \star \cdots \star (\alpha_{|\Pi|} \times \pi_{|\Pi|}),$ where α_j (1 \leq $j \leq$ |Π|) is a weighting constant.

The vector **P** is analogous to the vector π_j discussed in Section 2. Parallel to the discussion in Section 2, one needs to identify the set of boundary values of **P**. Let f_i be a function that maps **P** to ϕ_i . For the single system feature ϕ_i being considered, such a set is given by $\{P \mid (f_i(P) = \beta_i^{\max}) \bigvee (f_i(P) = \beta_i^{\min})\}.$

Let **P**orig be the assumed value of **P**. In addition, let $\mathbf{P}^{\star}(\phi_i)$ be, analogous to $\boldsymbol{\pi}_j^{\star}(\phi_i)$, the element in the set of boundary values such that the Euclidean distance from **P**^{orig} to $\mathbf{P}^{\star}(\phi_i)$, $\|\mathbf{P}^{\star}(\phi_i) - \mathbf{P}^{\text{orig}}\|_2$, is the smallest over all such distances from **P**orig to a point in the boundary set. Alternatively, the value $\|\mathbf{P}^{\star}(\phi_i) - \mathbf{P}^{\text{orig}}\|_2$ gives the largest Euclidean distance that the variable **P** can move in *any* direction from an assumed value of **P**orig without exceeding the tolerable limits on ϕ_i . Parallel to the discussion in Section 2, let the distance $\|\mathbf{P}^{\star}(\phi_i) - \mathbf{P}^{\text{orig}}\|_2$ be called the robustness radius, $r_{\mu}(\phi_i, \mathbf{P})$, of ϕ_i against

P. Mathematically,

$$
r_{\mu}(\phi_i, \mathbf{P}) = \min_{\mathbf{P}: (f_i(\mathbf{P}) = \beta_i^{\max}) \bigvee (f_i(\mathbf{P}) = \beta_i^{\min})} \|\mathbf{P} - \mathbf{P}^{\text{orig}}\|_2.
$$
\n(2)

Extending for all $\phi_i \in \Phi$, the robustness of resource allocation μ with respect to the performance feature set Φ against the perturbation parameter set Π is given by $\rho_{\mu}(\Phi, \mathbf{P}) = \min_{\phi_i \in \Phi} (r_{\mu}(\phi_i, \mathbf{P}))$.

So how can one use the value of robustness calculated in P-space? How can one relate this value to allowable changes in π_i 's? Consider a system that has perturbation parameters of different kinds. Assume that its robustness value, as calculated in P-space, is $r_{\mu}(\phi_i, \mathbf{P})$ for a set of π_j^{orig} values. To find out whether the system can operate without a constraint violation under a given set of π_j values, one can (a) convert the π_j values into a **P** value using the α_j 's, (b) calculate the distance of **P** from \mathbf{P}^{orig} as $\|\mathbf{P} - \mathbf{P}^{\text{orig}}\|_2$, and (c) determine if $\|\mathbf{P} - \mathbf{P}^{\text{orig}}\|_2 < r_{\mu}(\phi_i, \mathbf{P})$. If yes, then the system will not violate a constraint when operated at the given values of π_i 's.

The sensitivity-based weighting procedure for the calculation of α_i 's is now discussed. Typically, π_1 , π_2, \dots, π_{Π} will have different dimensions, i.e., will be measured in different units, e.g., seconds, objects per data set, bytes, etc. Before the concatenation of these vectors into **P**, they should be converted into a single dimension. The proposed preliminary approach in [2] suggested a sensitivity-based weighting, that is: α_j = $1/r_{\mu}(\phi_i, \pi_j)$. With this definition of α_j ,

$$
\mathbf{P} = \frac{\pi_1}{r_\mu(\phi_i, \pi_1)} \star \frac{\pi_2}{r_\mu(\phi_i, \pi_2)} \star \cdots \star \frac{\pi_{|\Pi|}}{r_\mu(\phi_i, \pi_{|\Pi|})}.
$$

Note that the units of $r_{\mu}(\phi_i, \pi_j)$ are the units of π_j . This fact renders **P** dimensionless. Therefore the robustness measurement, $r_{\mu}(\phi_i, \mathbf{P})$, would be dimensionless as well. However, the investigation done in this research shows that the sensitivity-based weighting method has the following problem. When the performance feature is a linear function of one-element perturbation parameters, the robustness radius depends only on the number of perturbation parameters. Any change in the coefficients of ϕ_i function, or the original values of the perturbation parameters does not affect the robustness radius. This means that if two systems have performance features that are linear functions of the same number of perturbation parameters, then they will have the same robustness radius. With such characteristic in a robustness measure, it is impossible to compare the robustness of different systems that meet the criterion given above.

The above point is now illustrated in more detail. Let $\phi_i(\pi_1, \pi_2, \cdots, \pi_n)$ be a function of *n* perturbation vectors, $\pi_1, \pi_2, \dots, \pi_n$, of different kinds. Assume that

vectors $\pi_1, \pi_2, \dots, \pi_n$ have only one element each, i.e., $\pi_1 = \pi_{11} = \pi_1, \pi_2 = \pi_{21} = \pi_2, \dots, \pi_n = \pi_{n1}$ π_n . The discussion below will consider only β_i^{\max} as a constraint.

As a general linear case, let $\phi_i(\pi_1, \pi_2, \cdots, \pi_n)$ = $k_1\pi_1$ + $k_2\pi_2$ + \cdots + $k_n\pi_n$, and the original values of the perturbation parameters be $\pi_1^{\text{orig}}, \pi_2^{\text{orig}}, \cdots, \pi_n^{\text{orig}}$, respectively, and $\hat{\beta}_i^{\text{max}} = \beta \phi_i^{\text{orig}} =$ $\beta(k_1\pi_1^{\text{orig}}+k_2\pi_2^{\text{orig}}+\cdots+k_n\pi_n^{\text{orig}})$, where β is an arbitrary constant greater than 1. Note that this implies that β_i^{\max} is a function of original value of the performance feature, ϕ_i . This assumption does make sense because in many cases we limit the changes in ϕ_i to some percentage of its original value (e.g. in a Grid-like system, makespan should not exceed 1.2 times its original value). The following steps are performed to calculate the robustness radius.

Step 1: Determine the robustness radius with respect to π_j , $r_\mu(\phi_i, \pi_j)$ by setting π_m , $m \neq j$, to π_m^{orig} in the ϕ_i function. Then let α_j be $1/r_\mu(\phi_i, \pi_j)$:

Example for
$$
\pi_1
$$

\n $\pi_m = \pi_m^{\text{orig}}, \quad m \neq 1$:
\nThe constraint equation is $\phi_i|_{\pi_m = \pi_m^{\text{orig}}} = \beta_i^{\text{max}},$
\n $\eta_i^{\text{max}} = \beta_i^{\text{max}}$

or

$$
k_1\pi_1 + k_2\pi_2^{\text{orig}} + \cdots + k_n\pi_n^{\text{orig}} = \beta(k_1\pi_1^{\text{orig}} + k_2\pi_2^{\text{orig}} + \cdots + k_n\pi_n^{\text{orig}}).
$$

Solving the above equation for π_1 ,

$$
\pi_1 = \frac{\beta k_1 \pi_1^{\text{orig}} + (\beta - 1)(k_2 \pi_2^{\text{orig}} + \dots + k_n \pi_n^{\text{orig}})}{k_1}.
$$

Now the robustness radius can be calculated by applying Equation 1. Since it is a one-dimensional space, the minimum Euclidean distance between $\hat{\pi}_1^{\text{orig}}$ and the constraint curve is just the distance between π_1^{orig} and the point π_1 which was calculated above. That is $r_{\mu}(\phi_i, \pi_1) = \pi_1 - \pi_1^{\text{orig}}$, therefore
 $r_{\mu}(\phi_i, \pi_1) = \frac{\beta - 1}{k_1} (k_1 \pi_1^{\text{orig}} + k_2 \pi_2^{\text{orig}} + \dots + k_n \pi_n^{\text{orig}})$

$$
\alpha_1 = 1/r_\mu(\phi_i, \pi_1)
$$

=
$$
\frac{k_1}{(\beta - 1)(k_1 \pi_1^{\text{orig}} + k_2 \pi_2^{\text{orig}} + \dots + k_n \pi_n^{\text{orig}})}.
$$

(3)

For the general case, $\pi_m = \pi_m^{\text{orig}}$, $m \neq j$:

The constraint equation is $\phi_i|_{\pi_m = \pi_m^{\text{orig}}} = \beta_i^{\max}$, $m \neq j$ or

$$
\sum_{\substack{m=1 \ m \neq j}}^{n} k_m \pi_m^{\text{orig}} + k_j \pi_j = \beta \sum_{m=1}^{n} k_m \pi_m^{\text{orig}}.
$$

Solving the above equation for π_i ,

$$
\pi_j = \frac{\beta k_j \pi_j^{\text{orig}} + (\beta - 1) \sum_{m=1, m \neq j}^{n} k_m \pi_m^{\text{orig}}}{k_j}
$$

As before, since it is a one-dimensional space, the minimum Euclidean distance between π_j^{orig} and the constraint curve is just the distance between π_j^{orig} and the point π_j which was calculated above. That is $r_{\mu}(\phi_i, \pi_j) = \pi_j - \pi_j^{\text{orig}}$, then

$$
r_{\mu}(\phi_i, \pi_j) = \frac{\beta - 1}{k_j} \sum_{m=1}^n k_m \pi_m^{\text{orig}}
$$

$$
\alpha_j = 1/r_\mu(\phi_i, \pi_j)
$$

$$
= \frac{k_j}{(\beta - 1) \sum_{m=1}^n k_m \pi_m^{\text{orig}}}
$$

.

Step 2: Using the α_j 's obtained in Step 1, determine the relationships between corresponding elements of **P** and π_j 's. Deduce the function f_i as a function of **P** from these relationships, and then generate a robustness metric in P-space.

 $\mathbf{P} = [\alpha_1 \pi_1 \ \alpha_2 \pi_2 \ \cdots \ \alpha_n \pi_n].$

The constraint equation is $\phi_i = \beta_i^{\max}$, where the lefthand side would be simplified by replacing π_i 's with their equivalents in P_i 's,

$$
\phi_i = k_1 \pi_1 + k_2 \pi_2 + \dots + k_n \pi_n \n= k_1 P_1/\alpha_1 + k_2 P_2/\alpha_2 + \dots + k_n P_n/\alpha_n.
$$

Given that

$$
\alpha_j = \frac{k_j}{(\beta - 1) \sum_{m=1}^n k_m \pi_m^{\text{orig}}},
$$

the ϕ_i equation would be simplified as:

We will get:

$$
\phi_i = k_1 P_1 \frac{\binom{\beta - 1}{m} \sum_{m=1}^n k_m \pi_m^{\text{orig}}}{k_1} + \cdots
$$

$$
\frac{(\beta - 1) \sum_{m=1}^n k_m \pi_m^{\text{orig}}}{k_n}
$$

$$
= (P_1 + P_2 + \cdots + P_n)(\beta - 1) \sum_{m=1}^n k_m \pi_m^{\text{orig}}.
$$

The right-hand side of the constraint equation, β_i^{\max} , would be:

$$
\beta_i^{\max} = \beta \phi_i^{\text{orig}} \n= \beta (k_1 \pi_1^{\text{orig}} + k_2 \pi_2^{\text{orig}} + \dots + k_n \pi_n^{\text{orig}}) \n= \beta \sum_{m=1}^n k_m \pi_m^{\text{orig}}.
$$

Therefore the constraint equation, $\phi_i = \beta_i^{\text{max}}$, would be simplified as:

$$
P_1 + P_2 + \cdots + P_n = \beta/(\beta - 1).
$$

Then the robustness radius of ϕ_i with respect to **P**, that is the minimum Euclidean distance from \mathbf{P}^{orig} to the *n*-dimensional plane above. Recall that given an n -dimensional plane, $a_1x_1 + a_2x_2 + \cdots + a_nx_n = b$, the minimum distance between a point X_0 : $(x_{01}, x_{02}, \dots, x_{0n})$ and the plane is would be:

$$
d = \frac{|a_1x_{01} + a_2x_{02} + \dots + a_nx_{0n} - b|}{\sqrt{a_1^2 + a_2^2 + \dots + a_n^2}}.
$$
 (4)

Then $r_{\mu}(\phi_i, \mathbf{P})$ would be:

$$
r_{\mu}(\phi_i, \mathbf{P}) = \frac{|P_1^{\text{orig}} + P_2^{\text{orig}} + \dots + P_n^{\text{orig}} - \beta/(\beta - 1)|}{\sqrt{n}}.
$$

The above equation can be furthur simplified by replacing P_j^{orig} 's by their equivalents, $\alpha_j \pi_j^{\text{orig}}$, where

$$
\alpha_j \pi_j^{\text{orig}} = \frac{k_j \pi_j^{\text{orig}}}{(\beta - 1) \sum_{m=1}^n k_m \pi_m^{\text{orig}}},
$$

$$
r_{\mu}(\phi_i, \mathbf{P}) = \frac{\frac{\sum_{j=1}^{n} k_j \pi_j^{\text{orig}}}{(\beta - 1) \sum_{m=1}^{n} k_m \pi_m^{\text{orig}}} - \frac{\beta}{\beta - 1}}{\sqrt{n}} = \frac{\frac{1}{\beta - 1} - \frac{\beta}{\beta - 1}}{\sqrt{n}} = \frac{1}{\sqrt{n}}.
$$

One can see that regardless of the values of k_j 's, β and the original values of π_j 's, the robustness radius is equal to $\frac{1}{\sqrt{n}}$. The fact that an increase in the robustness requirement, β_i^{\max} , does not change the robustness value is troubling. This may not be a desired characteristic of a robustness measure. Therefore the sensitivity-based concatenation may not give a satisfactory robustness measure.

3.2. Normalizing the Robustness Measure with respect to the Original Values of Perturbation Parameters

The purpose of this section is to propose a method for robustness measure calculation which can be used for multiple kinds of perturbations, and be able to compare different cases. We still have to merge π_j 's into **P** so as to compose a dimensionless perturbation parameter vector, and then determine the robustness metric by taking the minimum over the robustness radii of all $\phi_i \in \Phi$, as described in the Overview section.

Let the vector π_j have n_{π_j} elements $(1 \leq j \leq |\Pi|)$, that is $\pi_j = [\pi_{j1} \quad \pi_{j2} \quad \cdots \quad \pi_{jn_{\pi}}].$ Also redefine **P** as a concatenation of vectors $\pi_1, \pi_2, \cdots, \pi_{|\Pi|}$ normalized by the original values of their elements:

$$
\mathbf{P} = [\pi_{11}/\pi_{11}^{\text{orig}} \ \pi_{12}/\pi_{12}^{\text{orig}} \ \cdots \ \pi_{1n_{\pi_1}}/\pi_{1n_{\pi_1}}^{\text{orig}}] \star \cdots \ \star [\pi_{|\Pi|1}/\pi_{|\Pi|1}^{\text{orig}} \ \pi_{|\Pi|2}/\pi_{|\Pi|2}^{\text{orig}} \ \cdots \ \pi_{|\Pi|n_{\pi_{|\Pi|}}}/\pi_{|\Pi|n_{\pi_{|\Pi|}}}^{\text{orig}}].
$$
\n(5)

Now \mathbf{P}^{orig} will always be $\begin{bmatrix} 1 & 1 & \cdots & 1 \end{bmatrix}$. Using this definition, the elements of **P** would be dimensionless, therefore $r_{\mu}(\phi_i, \mathbf{P})$ would be dimensionless as well. The

quantity $\rho_{\mu}(\Phi, \mathbf{P})$ is defined as the robustness of resource allocation μ with respect to the performance feature set Φ against the perturbation parameter set Π , and is given by the minimum of all robustness radii, or $\rho_\mu(\Phi, \mathbf{P}) = \min_{\phi_i \in \Phi} (r_\mu(\phi_i, \mathbf{P})).$

Furthermore, in this definition, the relative, not the absolute, changes of each perturbation parameter element with respect to its original value are examined. We believe that this method can be applied on systems with a single perturbation parameter as well as the systems with multiple perturbation parameters, to achieve a uniform dimensionless robustness measure. And this method does not have the problem present in [2] which was discussed in Section 3.1.

For the same general linear case which was discussed in Section 3.1, $r_{\mu}(\phi_i, \mathbf{P})$ is calculated as following. The concatenated vector, **P**, is derived from Equation (4) :

 $\mathbf{P} = [\pi_1/\pi_1^{\text{orig}} \ \pi_2/\pi_2^{\text{orig}} \ \cdots \ \pi_n/\pi_n^{\text{orig}}].$

The left-hand side of the constraint equation, ϕ_i = β_i^{max} , would be simplified by replacing π_j 's with their equivalents in P_j 's,

$$
\phi_i = k_1 \pi_1 + k_2 \pi_2 + \dots + k_n \pi_n \n= k_1 \pi_1^{\text{orig}} P_1 + k_2 \pi_2^{\text{orig}} P_2 + \dots + k_n \pi_n^{\text{orig}} P_n.
$$

The right-hand side of the constraint equation, β_i^{\max} , would be:

$$
\beta_i^{\max} = \beta \phi_i^{\text{orig}} \n= \beta (k_1 \pi_1^{\text{orig}} + k_2 \pi_2^{\text{orig}} + \dots + k_n \pi_n^{\text{orig}}) \n= \beta \sum_{m=1}^n k_m \pi_m^{\text{orig}}.
$$

Therefore the constraint equation, $\phi_i = \beta_i^{\text{max}}$, would be simplified as:

$$
\sum_{j=1}^{n} k_j \pi_j^{\text{orig}} P_j = \beta \sum_{m=1}^{n} k_m \pi_m^{\text{orig}}
$$

\n
$$
k_1 \pi_1^{\text{orig}} P_1 + k_2 \pi_2^{\text{orig}} P_2 + \dots + k_n \pi_n^{\text{orig}} P_n =
$$

\n
$$
\beta \sum_{m=1}^{n} k_m \pi_m^{\text{orig}}.
$$

Then the robustness radius of ϕ_i with respect to **P**, that is the minimum Euclidean distance from **P**orig to the n -dimensional plane above. Using Equation 4, the robustness radius would be:

$$
r_{\mu}(\phi_i, \mathbf{P}) = \frac{\left| \sum_{j=1}^{n} k_j \pi_j^{\text{orig}} P_j^{\text{orig}} - \beta \sum_{m=1}^{n} k_m \pi_m^{\text{orig}} \right|}{\sqrt{\sum_{m=1}^{n} (k_m \pi_m^{\text{orig}})^2}}
$$

Recalling that $P = \begin{bmatrix} 1 & 1 & \cdots & 1 \end{bmatrix}$, the above equation is simplified to:

$$
r_{\mu}(\phi_i, \mathbf{P}) = \frac{\left| \sum_{j=1}^n k_j \pi_j^{\text{orig}} - \beta \sum_{m=1}^n k_m \pi_m^{\text{orig}} \right|}{\sqrt{\sum_{m=1}^n (k_m \pi_m^{\text{orig}})^2}}
$$

$$
= (\beta - 1) \frac{\left| \sum_{j=1}^n k_j \pi_j^{\text{orig}} \right|}{\sqrt{\sum_{m=1}^n (k_m \pi_m^{\text{orig}})^2}}.
$$

One can see that now the robustness radius depends, as it should, on the values of k_j 's, β , and the original values of π_j 's.

4. Conclusions

Parallel and distributed heterogeneous computing systems may operate in environments that undergo unpredicted changes in environment or component parts of system. A measurement method for the robustness of such systems has been presented and developed in [2]. However, that approach was mainly focused on systems in which all perturbation parameters were of the same kind. In our paper, we focus on the multiple perturbation parameter case and on merging the different perturbation parameters into a single perturbation parameter. The investigation done in this research further strengthens the sensitivity-based weighting method presented in [2]. When the performance feature is a linear function of one-element perturbation parameters, the robustness radius given in [2] depends only on the number of perturbation parameters. In other words, regardless of the values of the robustness requirement, the coefficients of ϕ_i function, or the original values of the perturbation parameters, the robustness radius is equal to $\frac{1}{\sqrt{n}}$, where *n* is the number of perturbation parameters. This means that if two systems have performance features that are linear functions of the same number of perturbation parameters, then they will have the same robustness radius. We propose an alternative robustness metric formulation that avoids the problem mentioned above.

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Biographies

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Φ	the set of all performance features
ϕ_i	the <i>i</i> -th element in Φ
$\langle \beta_i^{\text{min}}, \beta_i^{\text{max}} \rangle$	a tuple that gives the bounds of the tolerable variation in ϕ_i
Π	the set of all perturbation parameters
π_i	the <i>j</i> -th element in Π
f_{ij}	the function that maps π_i to ϕ_i
$n_{\boldsymbol{\pi}_i}$	the dimension of vector π_i
μ	a resource allocation
$r_{\mu}(\phi_i, \pi_j)$	the robustness radius of resource allocation μ with respect to ϕ_i against π_i
$\rho_\mu(\Phi, \pi_i)$	the robustness of resource allocation μ with respect to set Φ against π_i
f_i	the function that maps P to ϕ_i
P	a weighted concatenation of the vectors $\pi_1, \pi_2, \cdots, \pi_{ \Pi }$
$r_{\mu}(\phi_i, \mathbf{P})$	the robustness radius of resource allocation μ with respect to ϕ_i against P
$\rho_\mu(\Phi, \mathbf{P})$	the robustness of resource allocation μ with respect to set Φ against P

Table 1. Glossary of notation.

Figure 1. Some possible directions of increase of the perturbation parameter π_j , and the direc- ${\bf t}$ ion of the smallest increase. The curve plots the set of points, $\{\pi_j | \: f_{ij}(\pi_j) = \beta_i^{\max}\}$. The set of boundary points, $\{\bm\pi_j|~ f_{ij}(\bm\pi_j)=\beta^{\min}_i\}$ is given by the points on the π_{j1} -axis and π_{j2} -axis.

